

Digital electronics

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March 17, 2021

Plan

7 lessons/lab sessions (4-hour long schedules):

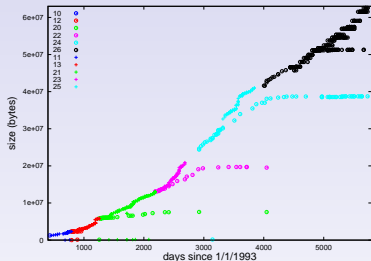
- 1 Executive environments: principles and introduction, getting started with FreeRTOS
- 2 FreeRTOS, RTEMS, Nuttx ... multitasking and associated methods to make sure shared data and resources are kept in known states (mutex & semaphore)
- 3 Using the scheduler, mutex and semaphores to solve the “philosopher problem”
- 4 Embedded systems with GNU/Linux – POSIX compatible operating system
Architecture of an operating system, kernel v.s userspace
Internet connectivity and networking
- 5 Accessing hardware resources from userspace – memory translation from physical to virtual address space (Memory Management Unit) – /dev/mem
- 6 Accessing hardware resources from a web server – internet connected instrument
- 7 From userspace to kernel space: character device (*char device*) for communicating between users and the kernel

Introduction: kernel & kernel modules

- Developer \Rightarrow focus on the software-hardware interface
- Only the kernel can access all resources of the processor (e.g. interrupts, DMA)
- The kernel shares resources and make sure hardware access is consistent (kernel = supervisor)
- Adding new functionalities to the kernel: modules (“plugins”)

Porting GNU/Linux to a new architecture

- a (cross-)compilation toolchain
- define the memory map of the target (-T option of ld)
- bootloader development (peripheral initialization and loading the kernel to memory)¹
- developing kernel modules for handling hardware dedicated to the new architecture².

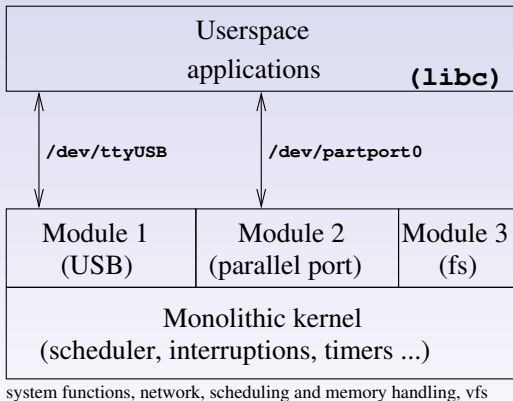


Evolution of the Linux kernel
.tar.gz archive size (www.kernel.org)

¹S. Guinot, J.-M Friedt, *GNU/Linux sur Playstation Portable*, GNU/Linux Magazine France **114**, Mars 2009, pp.30-40

²J. Corbet, A. Rubini & G. Kroah-Hartman, *Linux Device Drivers, 3rd Edition*, O'Reilly (2005), available at <http://lwn.net/Kernel/LDD3/>

Operating System (OS) architecture



The /dev directory

- Pseudo-files linking userspace and kernelspace through system calls implemented in `libc` (`man syscalls`)
- each *device* is defined with a *major number* and, within each class, a *minor number*
- the kernel is not aware of filesystems: on the kernel side, a driver is associated with a given major number and handles each syscall:

```
static struct file_operations qadc_fops =
{
    owner:           THIS_MODULE,
    read:            qadc_read,
    open:            qadc_open,
    ioctl:           qadc_ioctl,
    release:         qadc_release,
};

static int __init qadc_init (void)
{...
    register_chrdev (qadc_major, "ppp", &qadc_fops);
    ...}

module_init (qadc_init);
module_exit (qadc_exit);
```

Basic structure

Kernel module \Rightarrow no `main()` starting function:

- one entry and one exit function:
`module_init (my_init_function);` and
`module_exit (my_exit_function);`
- initializing the communication link with userspace:
`register_chrdev (90 , "jmf" ,&fops);`
- information displayed through `/var/log/syslog` using
`printk (KERN_ALERT "formatted message");` (*no coma !*)
- ability to dynamically create dev entries using *misc. devices*

```
int hello_start() // init_module(void)
{
    jmfdev.name = "jmf"; // /dev/jmf: major = classe misc
    jmfdev.minor = MISC_DYNAMIC_MINOR;
    jmfdev.fops = &fops;
    misc_register(&jmfdev); // dynamic creation of /dev/jmf
}
```

Module handling commands

Modules (*kernel object*) with `.ko` extension

- are linked to the kernel with `insmod toto.ko`,
- are removed from the kernel with `rmmmod toto`.
- list of attached modules with `lsmod`
- Messages logged by modules using `printk` are displayed with `dmesg`
- “Official” modules are located in `/lib/modules`
- “Official” modules with dependencies are loaded with `modprobe`³

³see `/lib/modules/version/modules.dep`

Directories needed for compiling modules

The sources of the kernel running on the PC are located in `/usr/src`:

- Debian/GNU Linux separates generic headers from parts specific to a given architecture
- `/usr/src/linux-headers-*-common` does *not* include parts specific to the current architecture
- `/usr/src/linux-headers-*-*arch` is the directory we use to find the specific configuration of the kernel running on the host.

New Makefile architecture for compiling modules relying on the kernel Makefile by adding the new kernel target to the `obj-m` variable with

```
obj-m +=0mymod.o 1mymod.o 2mymod.o
```

```
all:
```

```
    make -C /usr/src/linux-headers-4.8.0-1-amd64 M=$(PWD) modules
```

For cross-compiling to the target architecture: `output/build/linux-*` in Buildroot provides the Linux kernel

```
obj-m +=0mymod.o 1mymod.o 2mymod.o
```

```
all:
```

```
    make ARCH=arm CROSS_COMPILE=arm-linux- -C \  
    ../buildroot-2020.11.1/output/build/linux-custom/ M=$(PWD) modules
```

Real v.s virtual addresses

Same problem as `/dev/mem` and `mmap` in userspace: here too the MMU must provide the virtual block address translated to the physical address:

```
request_mem_region (IO_BASE, 36*9, "GPIO test" );// PG9
jmf_gpio = ( u32 )ioremap(IO_BASE, 36*9);      // 36 bytes/port, ports A-I
writel (0x10, jmf_gpio+36*6+0x04 );           // GPIO PG9 en sortie
```

30.2. Port Register List

Module Name	Base Address	
PIO	0x01C20800	
Register Name	Offset	Description
Pn_CFG0	n*0x24+0x00	Port n Configure Register 0 (n from 0 to 9)
Pn_CFG1	n*0x24+0x04	Port n Configure Register 1 (n from 0 to 9)
Pn_CFG2	n*0x24+0x08	Port n Configure Register 2 (n from 0 to 9)
Pn_CFG3	n*0x24+0x0C	Port n Configure Register 3 (n from 0 to 9)
Pn_DAT	n*0x24+0x10	Port n Data Register (n from 0 to 9)
Pn_DRV0	n*0x24+0x14	Port n Multi-Driving Register 0 (n from 0 to 9)
Pn_DRV1	n*0x24+0x18	Port n Multi-Driving Register 1 (n from 0 to 9)
Pn_PUL0	n*0x24+0x1C	Port n Pull Register 0 (n from 0 to 9)

Addresses given in the microprocessor datasheet (here Allwinner A10)

Methods provided by the kernel

Many functionalities provided by the kernel for handling tasks (scheduler), timer or communicating with userspace.

- concepts close to microcontroller programming

```
init_timer_on_stack (&exp_timer );  
exp_timer.expires = jiffies+delay*HZ ;  
exp_timer.data = 0;  
exp_timer.function = do_something ;  
add_timer (&exp_timer);
```

- *tasklet*: software interrupts

```
char my_tasklet_data[]="my_tasklet_function was called";  
void my_tasklet_function( unsigned long data )  
{ printk( "%s\n", (char *)data ); return; }
```

```
DECLARE_TASKLET( my_tasklet, my_tasklet_function,  
                (unsigned long) &my_tasklet_data );
```

```
int init_module( void )  
{ tasklet_schedule( &my_tasklet ); return 0; }
```

```
void cleanup_module( void )  
{ tasklet_kill( &my_tasklet ); return; }
```

→ ISR should be **short** and **schedule** a tasklet

Kernel module: interrupt handling

Only a kernel module can handle interrupts (e.g. interrupt 7 is associated with pin 10 of the parallel port on a PC)

```
jmfriedt@(none):~$ cat /proc/interrupts | grep parp
7:          0      IO-APIC-edge  parport0
```

handled with ⁴

```
#define PARALLEL_PORT_INTERRUPT 7
#define BASEPORT                0x378
int s=0;
static int int_handler(void) {printk(">>> PARALLEL PORT INT\n");s=1;}
static int __init jmfpport_init_module (void) {
    ret = request_irq(PARALLEL_PORT_INTERRUPT, &int_handler,
                     SA_INTERRUPT, "parallelport", NULL);

    enable_irq(7);
    outb_p(0x10, BASEPORT + 2);
}
static ssize_t skeleton_read (struct file *file, char *buf,
                              size_t count, loff_t *ppos) {
    [...]
    do {} while (sortir==0); // interruptible_sleep_on(&skeleton_wait);
    err = copy_to_user(buf, string, counter);
    ...
}
```

⁴www.captain.at/howto-linux-device-driver-template-skeleton.php

Practical demonstration

- ① Kernel module for Linux
- ② Converting from virtual to real memory in kernel space
- ③ Dynamic communication entry creation in `/dev`
- ④ Reading a datasheet (hardware addresses) remains mandatory !
- ⑤ GNU License to be allowed to access the functionalities provided by the kernel or other modules (hardware abstraction)

→ convert the userspace C blinking LED example to a kernel module